# Schlüsselgerät 41 

Technical aspects of the German WWII Hitlermühle
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## Contents

Abstract ..... 3
Introduction ..... 5
Pseudo-random number generator ..... 6
Operation ..... 9
Internal Key ..... 9
External Key ..... 9
Ciphering and deciphering ..... 9
Pin-wheels ..... 11
Print head ..... 13
Encryption ..... 13
Decryption ..... 13
Wheel stepping ..... 15
Hagelin pin-and-lug machines until 1943 ..... 15
Schlüsselgerät 41 ..... 15
Order of events ..... 16
Creation of the PRN ..... 18
Horizontal movement of the bars ..... 19
Sensing arms ..... 20
Inactive pin ..... 21
Active pin ..... 21
Inversion ..... 21
Simulation ..... 22
Conclusions ..... 24
About the authors ..... 25
References ..... 25
Links ..... 24

## Figures

1 IPO-model ..... 5
2 Keyboard layout ..... 5
3 Pin-wheel seen from the left and front ..... 6
4 The six pin-wheels and the yellow release button ..... 6
4 Overview of the controls and features ..... 7
5 Left side of the SG-41 ..... 8
6 Right side of the SG-41 ..... 8
7 Sensing position on the pin-wheel (seen from the left side) ..... 11
8 Example of a bitmask of the first five wheels ..... 11
9 Example of an inverted bitmask ..... 11
10 Print head during encryption (left) and decryption (right) ..... 13
11 Sensing positions on the pin-wheels (seen from the left side) ..... 15
12 Order of events represented as the hours of a clock ..... 17
13 Bar with a lug at the position of pin-wheel 1 ..... 18
14 Top view of the cage with the sideways movable bars ..... 18
15 Value of each wheel (bitmask) ..... 19
16 Sliding bar shown at different positions whilst the cage revolves ..... 19
17 Three possible positions of a bar ..... 20
18 Sickle-shaped sensing arms ..... 20
19 Behavior when wheel $6=$ ' 0 ' (left) and when wheel $6=$ ' 1 ' (right) ..... 21
20 HELL H-54 (left) and Hagelin CX-52 (right) ..... 22
Tables
1 Layout of the pin-wheels ..... 10
2 Layout of the print head ..... 12
3 Examples of encryption and decryption ..... 12
4 Chronological order of events ..... 16
5 Pin-settings (example) ..... 22
6 Wheel positions when encrypting the given plaintext (PT) into ciphertext (CT) ..... 23

## Abstract

This paper describes the exact operation of Schlüsselgerät 41 (cipher machine 41), also known as SG-41 or Hitlermüble (Hitler mill), and explains the complex order and timing of the events that take place when a text is enciphered. The authors believe that this is the first time that some of the properties of this machine are revealed. A comparison is made with the contemporary and similar Hagelin M-209 and BC-38 machines, and it is shown that SG-41 was more advanced in several ways. Furthermore, the authors present the wheel positions for a finite number of steps, along with the plaintext, the ciphertext and the required settings. The strength of the cipher and possible cryptanalytic attacks will not be discussed here. This paper is largely the result of the work of Klaus Kopacz, carried out whilst restoring a broken SG-41, whilst Paul Reuvers is responsible for the English translation, the drawings, additional descriptions and some of the photographs.

## Introduction

Schlüsselgerät 41 (German: cipher machine 41), commonly known as SG-41 or the Hitlermüble (German: Hitler mill), was a mechanical cipher machine, developed and built during World War II (WWII) by Wanderer Werke in Siegmar Schönau, known today as Chemnitz (Germany). From 1943 to 1945 it was used by the German war machine - in particular by the German Intelligence Service, the Abwehr. The SG-41 is a so-called pin-and-lug cipher machine, and shows great resemblance to the Hagelin cipher machines of the era. Nevertheless, there are significant differences.

There are two versions of the machine: (1) the basic SG-41, with 26 keys for the Latin alphabet (A-Z) and (2) the SG-41Z, with 10 numbered keys ( $0-9$ ), suitable for numeric messages only. The SG-41Z was used by the Wetterdienst (weather service) of the Luftwaffe (German Air Force), and will be described in a separate article. The current paper only describes the basic SG-41.


Figure 1 - IPO-model
Like all processing devices, the SG-41 consists of three basic elements: (1) a keyboard for the input, (2) a mechanical pseudorandom number generator (PRNG) for processing the text (the actual encryption or 'cipher'), and (3) a printer for the output. The keyboard has 26 keys that are marked with the uppercase letters of the Latin alphabet (A-Z) in the German layout. All letters are black on a white key top, except for the letter 'J', which has a red key top. In the plaintext it is automatically replaced by a space (blank). As the letter 'J' cannot be used in the plaintext, it has to be substituted by the letter 'I'. Behind the upper row of keys is a label with the numerals $1-9$ and 0 . It allows the keys of the upper row ( $\mathrm{Q} \ldots \mathrm{P}$ ) to be used for numbers after the user has entered a predefined shift-sequence. To switch back to letters, the user has to enter a predefined unshift-sequence.


Figure 2 - Keyboard layout
When encrypting or decrypting, one of the print heads is first rotated to the selected letter. The letter is then printed onto a paper strip, after which all print heads are rotated an arbitrary number of steps and the enciphered letter is printed on another paper strip. The arbitrary number of steps is determined by the mechanical pseudo-random number generator (PRNG).

After entering a letter on the keyboard, the user has to rotate the hand crank at the right side of the machine clockwise by one full revolution. This causes the letter to be encrypted or decrypted, depending on the selected mode of operation. At the end of the revolution, the plaintext and ciphertext letters have been printed onto individual paper strips, and the PRNG is set to its next state. At the same time, the output letter is shown in a window.

## Pseudo-random number generator

The PRNG uses six pin-wheels, each with a different number of segments, plus a rotating drum, or cage, with horizontally movable slide bars that multiply the value of the current pin-wheel with the presence of an active pin, and adds it to a running total. Each pin-wheel has a series of pins on its circumference, equal to the number of segments on the wheel. Each pin can be shifted to the left (inactive) or to the right (active). These states are also known as ' 0 ' (inactive) and ' 1 ' (active). The state of the pins can be programmed in situ. Resetting all pins to the inactive state (0), requires the F/L switch to be set to the 'L' position, the key 'Löschen' to be held down, and the crank to make 25 full revolutions. This procedure will be described later.


Figure 3
Pinwheel seen from the left and front showing the inactive (left) and active (right) states

The Basic Setting of the 144 pins, also known as the Internal Key or Daily Key, allows $2^{144}$ different combinations, which is equivalent to approx. $2.2 \times 10^{43}$. In addition, the six pin-wheels must be set to their initial position before encrypting or decrypting a text. This is known as the External Key, or Message Key. There are $25 \times 25 \times 23 \times 23 \times 24 \times 24=190,440,000$ possible combinations for the External Key.

The pseudo-random number (PRN), which is used to encrypt or decrypt a letter, is built by the PRNG from the state of the pins at the circumference of the six pin-wheels, at a given position. In this process, the 6th wheel (i.e. the rightmost wheel) drives an inverting function, which was a novely as it had never been implemented in a comparable machine before. This function is also known as the complementary feature.

The stepping of the six pin-wheels evolves in four stages, driven by the state of the pins of all six pin-wheels, sensed at a specific position on the circumference of the wheels. The ideas of the pin-wheels, the cylindrical cage with sideways movable bars, and the circular print head, were clearly 'borrowed' from contemporary Hagelin machines like the M-209 and the BC-38. It should be noted however, that the wheel stepping and the generation of the PRN are significantly improved in the SG- 41 .


Figure 4 - The six pin-wheels and the yellow release button

The machine has a built-in double printer, that produces the output on two narrow paper strips, with the plaintext always on the right, and the ciphertext always on the left. This fixed assignment of the printers is accomplished by shifting the print head - which has four letter-rings - under control of the $\mathrm{V} / \mathrm{E}$ switch ( $\mathrm{V}=$ Verschlüsseln = encrypting, $\mathrm{E}=$ Entschlüsseln $=$ decrypting). This is an improvement over the Hagelin BC-38. The output of the ciphertext printer (left) is automatically formatted into five-letter groups, that are separated by a space. In addition, the paper of the leftmost printer can be advanced in fixed steps with a lever - marked 'Papier' (paper) - at the left side of the machine.

When loading or positioning the paper strips, the key marked 'Papier' (paper) on top of the machine can be pressed to release the mechanism, so that the paper strip can be moved freely. At the bottom of the machine is a drawer that holds the two paper supply spools: one for the plaintext and one for the ciphertext.

When sending a message, the sheet with the plaintext can be placed in a hinged copy holder on top of the machine. When entering a text, a resettable counter at the top surface of the machine counts the number of encrypted letters. Also at the top surface is a removable ink cartridge with four ink rolls, that can be inserted in two ways, marked ' 1 ' and ' 2 '. In each position, two of the ink rolls are placed before the print head. The other two are spares and become active when the cartridge is reversed.


Figure 5 - Overview of the controls and features


Figure 6 - Left side of the SG-41


Figure 7 - Right side of the SG-41

## Operation

## Internal Key

Each pin-wheel has a fixed number of segments, or positions. For each position there is a pin on the side of the wheel, that can be shifted to the left (inactive) or right (active). Before setting the Internal Key, also known as the Daily Key, all 144 pins must be reset (i.e. set to the left). During the reset procedure, the cage and the printer will be disabled temporarily. Resetting goes as follows:

- Set the F/L knob to the 'L' position ( $\mathrm{L}=$ Löschen $=$ clear ).
- Release the handle marked 'Löschen', shift it to the left and hold it in this position.
- Make at least 25 full revolutions with the crank until the display reads ' 025 '.
- Release the 'Löschen' handle.
- Reset the character counter by rotating its knob until the counter reads ' 000 '.
- Set the F/L knob back to the ' F ' position ( $\mathrm{F}=$ Funktion = operate) .

All 144 pins are now set to the inactive state (' 0 '). The active pins can now be set as follows:

- Open the hinged window over the pin-wheels.
- Push the yellow-marked release button to the right of the pin-wheels (see Figure 4).
- The wheels can now be rotated freely and the relevant pins can be set to the active state (right).

The wheels will automatically be locked again as soon as the crank is operated.

## External Key

Before encrypting or decrypting a text, the six pin-wheels must be set to the current External Key, also known as the Message Key. The segments of the leftmost four wheels are marked with letters, whilst the rightmost two have numbers.

- Place the crank in the lowest position (i.e. at 6 o'clock).
- Open the hinged window over the pin-wheels.
- Push the yellow-marked release button to the right of the pin-wheels.
- Turn each rotor so that the desired letter or number will be visible through the window.
- Close the window when ready.


## Ciphering and deciphering

Encryption and decryption of a message goes as follows:

- Set the F/L knob to 'F'.
- Push and hold the knob marked 'Einfärbung Papier'.
- Pull-out several centimeters of each of the paper strips, and tear them off.
- Release the knob.
- Set the V/E knob to 'V' for encryption, or 'E' for decryption.
- Reset the character counter.
- Enter a letter on the keyboard and turn the crank clockwise until it locks again in the 6 o'clock position.
- Repeat the above step until all letters of the message have been processed.

Each time a letter is entered and the crank is rotated, the plaintext letter and the ciphertext letter are printed. A space is inserted automatically after each five-letter group in the ciphertext only. If necessary, group spaces can be inserted manually by operating the 'Papier' lever (see Figure 6). When the message is complete, press the 'Papier' knob and pull out the paper strips as far as necessary, before tearing them off.

|  |  | 1 | 2 | 3 | 4 | 5 | 6 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 25 | 25 |  |  |  |  |
|  | 25 | Z | Z |  |  | 24 | 24 |
|  | 24 | Y | Y | 23 | 23 | 24 | 57 |
|  | 23 | X | X | X | X | 23 | 55 |
|  | 22 | W | W | W | W | 22 | 52 |
|  | 21 | V | V | V | V | 21 | 50 |
|  | 20 | U | U | U | U | 20 | 47 |
|  | 19 | T | T | T | T | 19 | 45 |
|  | 18 | S | S | S | S | 18 | 42 |
|  | 17 | R | R | R | R | 17 | 40 |
|  | 16 | Q | Q | Q | Q | 16 | 37 |
|  | 15 | P | P | P | P | 15 | 35 |
| Window <br> Sensed | 14 | O | 0 | 0 | O | 14 | 32 |
|  | 13 | N | N | N | N | 13 | 30 |
|  | 12 | M | M | M | M | 12 | 27 |
|  | 11 | L | L | L | L | 11 | 25 |
|  | 10 | K | K | K | K | 10 | 22 |
|  | 9 | I | I | 1 | 1 | 09 | 20 |
|  | 8 | H | H | H | H | 08 | 17 |
|  | 7 | G | G | G | G | 07 | 15 |
|  | 6 | F | F | F | F | 06 | 12 |
|  | 5 | E | E | E | E | 05 | 10 |
|  | 4 | D | D | D | D | 04 | 07 |
|  | 3 | C | C | C | C | 03 | 05 |
|  | 2 | B | B | B | B | 02 | 02 |
|  | 1 | A | A | A | A | 01 | 00 |
| ValuePin sensing offset |  | 1 | 2 | 4 | 8 | 10 | inv. |
|  |  | +8 | +8 | +8 | +8 | +8 | +8 |

Table 1 - Layout of the pin-wheels

## Pin-wheels

At the heart of the pseudo-random number generator (PRNG) are the six pin-wheels that are located below a hinged window at the top, numbered 1 to 6 from left to right. The wheels are fitted permanently on an axle and cannot be swapped. Each wheel has a different number of segments, or positions: $25,25,23,23,24$, and 24 respectively. Each segment is marked with a letter or number, in ascending order from bottom to top, as shown in Table 1. When stepping, the front face of a wheel moves downwards, so that the numbers and letters pass by the window in ascending order.


Figure 8 - Sensing position on the pin-wheel (seen from the left side)
When building the pseudo-random number (PRN), the state of a pin is sensed at a certain position on the wheel, which is different from the position shown in the window. This position is 8 steps above the letter that is shown in the window, as illustrated in Figure 8. A few examples that are highlighted in Table 1:

- When wheel 1 shows the letter ' W ' in the window (position 22 ), the state of the pin at position $(22+8)$ modulo $25=5$ is sensed, which is at the letter ' $E$ '.
- When wheel 5 shows the number 10 in the window, the state of the pin at position $(10+8)$ modulo $24=18$ is sensed.

When a sensed pin of wheels 1 to 5 is active, the value of that wheel is added to a running total. These values are $1,2,4,8$ and 10 respectively. In many respects this behaves like a binary bitmask, albeit with different values for the bits. For example: if the sensed pins on wheels 1,3 and 4 are active, the bitmask can be written as 10110 and the total value will be $1+4+8=13$. Likewise, if all pins are active (bitmask 11111), the total value will be $1+2+4+8+10=25$.


Figure 9 - Example of a bitmask of the first five wheels
When the sensed pin of wheel 6 is active, the bitmask will be inverted. For example: if the bitmask was 10110 and the sensed pin of wheel 6 is active (i.e. '1'), the bitmask becomes 01001 , which results in a total value of $2+10=12$.


Figure 10 - Example of an inverted bitmask

|  | Ciphertext |  | Plaintext |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | 4 |
| 26 | E | I | I | E |
| 25 | U | D | D | U |
| 24 | Y | Z | Z | Y |
| 23 | L | Q | Q | L |
| 22 | C | N | N | C |
| 21 | 0 | X | X | O |
| 20 | S | M | M | S |
| 19 | G | T | T | G |
| 18 | V | B | B | V |
| 17 | J | W | W | J |
| 16 | W | J | J | W |
| 15 | B | V | V | B |
| 14 | T | G | G | T |
| 13 | M | S | S | M |
| 12 | X | 0 | 0 | X |
| 11 | N | C | C | N |
| 10 | Q | L | L | Q |
| 9 | Z | Y | Y | Z |
| 8 | D | U | U | D |
| 7 | I | E | E | 1 |
| 6 | F | P | P | F |
| 5 | H | A | A | H |
| 4 | R | K | K | R |
| 3 | K | R | R | K |
| 2 | A | H | H | A |
| 1 | P | F | F | P |
|  |  |  |  |  |
|  |  |  |  |  |

Table 2 - Layout of the print head Input

Encryption

| 2 | 4 |  | 1 | 3 |
| :---: | :---: | :---: | :---: | :---: |
| 1 I | E |  | E | 1 |
| D | U |  | U | D |
| Z | Y |  | Y | Z |
| Q | L |  | L | 0 |
| N ${ }_{\text {a }}^{\text {a }}$ | C |  | C | N |
| X 둔 | 0 |  | 0 | X |
| M | S |  | S | M |
| T | G |  | G | T |
| B | V | nput | V | B |
| W | J |  | J | W |
| J | W |  | W | J |
| V | B |  | B | v |
| G | T |  | T | G |
| S | M |  | M | S |
| 0 | X |  | X | 0 |
| C | N |  | N | C |
| L | 0 |  | 0 | L |
| Y | Z |  | Z | Y |
| U | D |  | D | U |
| E | 1 | input $\rightarrow$ | 1 | E |
| P | F |  | F | P |
| A | H |  | H | A |
| K | R |  | R | K |
| R | K |  | K | R |
| H | A |  | A | H |
| F | P |  | P | F |

## Print head

The print head consists of four cylindrical rings, mounted on a single axle. The layout of the four rings - numbered 1 to 4 - is given in Table 2. The outer two rings ( 1 and 4 ) are identical, just like the inner two rings ( 2 and 3 ). When the print head is moved by the Pseudo-Random Number Generator (PRNG), the letters on the front face of the print head move downwards, and appear at the print position as defined in Table 2, reading the table from bottom to top.

Each letter ring has the 26 letters of the Latin alphabet in a fixed scrambled order. The order of the letters on the inner rings is the reverse of the order of the letters on the outer rings. In Table 2, three letters (PAK) are highlighted to illustrate the reversed order. It allows the machine to encrypt and decrypt without changing the inner mechanics. Had reverse letter rings not been used, the rotational direction of the print head had to be reversed during decryption. The print head is shown in Figure 11.


Figure 11 - Print head during encryption (left) and decryption (right)
The V/E knob - used to select between encryption (V) and decryption (E) - only shifts the print head sideways, so that the plaintext is always printed on the right paper strip, and the ciphertext is always printed on the left paper strip. A similar system - with three letter rings - was used in the Hagelin BC-38. Later Hagelin machines, such as the C-446 and CX-52, had only two letter rings on the print head, which could not be shifted sideways. On the Hagelin machines the C/D knob - used to select between encryption (C) and decryption (D) - enables the 5-letter group spacing only in encryption mode.

Examples for encrypting and decrypting a letter, are given in Table 3. In the examples, the PRNG is assumed to have produced the pseudo-random number (PRN) 8, which means that the print head is rotated by 8 positions as part of the encryption/ decryption process. The machine is assumed to be in the same state for both examples, which are further explained below.

## Encryption

The V/E knob is set to ' $V$ '. When encrypting the letter ' $V$ ', the user presses the V-key on the keyboard and starts turning the crank. The print head is turned to position 18 and the letter ' V ' is printed on the rightmost paper strip. Without the PRNG, the letter ' B ' would have been selected as the ciphertext. When continuing to turn the crank however, the print head will be turned by a number of positions - 8 in this case - as a result of the PRN. It then stops at position 26 and prints the letter 'I' on the leftmost paper strip.

## Decryption

The V/E knob is set to ' E '. When decrypting the letter ' I ', the user presses the I-key on the keyboard and starts turning the crank. The print head is turned to position 7 and the letter 'I' is printed on the leftmost paper strip. Without the PRNG, the letter ' E ' would have been selected as the plaintext. When continuing to turn the crank however, the print head will be turned by a number of positions - 8 in this case - as a result of the PRN. It then stops at position 15 and prints the letter ' V ' on the rightmost paper strip.

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## Wheel stepping

## Hagelin pin-and-lug cipher machines until 1943

Up to 1943, the well-known mechanical Hagelin pin-and-lug machines, also known as the C-machines, used a cylindrical drum with sideways movable bars at the heart of its PRNG. This drum is also known as the 'cage'. It was based on a money exchange machine that Boris Hagelin had developed in the early 1930s [1]. The function of the cage was the addition of different values based on a bitmask. For every active ' 1 ' bit in the bitmask, the value of that bit was added to a running total. In the C-machines the sum of the active bit values was used to generate a variable rotation of the print head.

Each pin-wheel has a different number of segments, or positions, and is rotated by one position each time a letter is encrypted. For each encrypted letter, the cage makes one full revolution ( $360^{\circ}$. Active pins of a pin-wheel are sensed by an internal lever, and causes a bar that carries a lug at that position, to shift from right to left, were it engages a cogwheel, which in turn rotates the print head by one position. Some C-machines had fixed lugs on the bars, whereas others had configurable ones. Each bar arrives at the cogwheel in the rightmost (inactive) or leftmost (active) position. All pin-wheels contribute with their sensed pins (in combination with their value) to the creation of the PRN, which is the number of positions the print head is advanced.

## Schlüsselgerät 41

Stepping of the pin-wheels of the SG-41 is controlled by the rotation of the crank in four phases (I, II, III and IV) and is based on the state of the pins at certain positions on the circumference of the pin-wheels. Two of the stepping phases are carried out before the encryption of a letter. The other two stepping phases are carried out after the encryption of a letter and will therefore have an effect on the next letter that is to be encrypted. This results in a complex stepping pattern of 1 to 4 positions for pinwheels 2 to 6 , and 1 or 2 positions for wheel 1 .


Figure 12 - Sensing positions on the pin-wheel (seen from the left lide)
The sensing position of the pins, used by the stepping mechanism, is not the same as the one that is used by the cage to calculate the rotation of the print head. For the stepping mechanism, the state of the pins 5 positions below the one that is shown in the window are sensed, as illustrated in Figure 12 (the blue arrow). These positions are latched before the first two stepping phases (I and II) are executed. Before the third stepping phase (III), the pins are re-sensed, which is after the wheels might have moved. In the fourth phase (IV), all wheels are stepped by one position.

## Order of events

After entering a letter, the crank has to make one full revolution $\left(360^{\circ}\right)$ to get the enciphered letter. The crank operates the internal mechanism, causes the PRNG to create a PRN, rotates the print head, (conditionally) steps the pin-wheels, and causes the input and output letters to be printed onto two paper strips. Several events take place during the $360^{\circ}$ revolution, at different positions of the crank. In order to explain the operation of the SG-41 and the order in which the events take place, the full revolution is divided into 12 segments, named after the hours of a clock, as illustrated in Figure 13. The same events are listed in chronological order in Table 4.

In rest, the crank is in the 6 o' clock position and is locked until a key is pressed. Once a key is pressed, the crank is released and must be turned $360^{\circ}$ clockwise to encrypt the letter, after which the crank is locked again. Immediately at the start, between 6 and 7 o'clock, the print head driving cylinder is wound up. It has several lugs at its circumference - one for each letter of the alphabet - and is visible in Figure 15. When the cylinder is released at 7 o'clock, it causes the print head to rotate until one of its lugs is stopped by the raised lever of the currently pressed key. At this point, the print head stops at the selected letter.


Table 4 - Chronological order of events

If pin of wheel 6 was ' 1 ' at 10 o'clock position, all pin wheels are stepped by one position


Figure 13 - Order of events represented as the hours of a clock

Stepping of the pin-wheels at 10-11 o'clock (Phase I) and 11-12 o'clock (Phase II) takes place before the PRN is generated and the print head is rotated. This means that they have an effect on the encryption of the current input letter. The last stepping phases, at and 4-5 o'clock (Phase III) and 5-6 o'clock (Phase IV), take place after the letter has been encrypted, and therefore only have an effect on the next letter that is to be encrypted.

Stepping of the individual wheels - e.g. at 10-11 o'clock (Phase I) - does not take place sequentially, but in parallel. The decision to step a wheel (or not) has been sensed and latched earlier and is carried out in parallel at this stage. In other words: if a pin of a given pin-wheel ( n ) is active (' 1 '), the wheel to its right ( $\mathrm{n}+1$ ) makes a single step, regardless of whether or not wheel $(\mathrm{n})$ was stepped itself by the wheel to its left $(\mathrm{n}-1)$.

For wheel stepping, the state of the pins of the pin-wheels is sensed at 5 positions under the window and memorised (latched) for the 10-11 o'clock (Phase I) and 11-12 o'clock (Phase II) stepping phases. Before the 4-5 o'clock stepping phase (III), the state of the pins is re-sensed.

Stepping at $10-11$ o'clock (Phase I) only takes place if the sensed pin of wheel 6 is active (' 1 '). It is possible that wheel 6 is stepped by wheel 5 (if wheel 5 had an active pin) but this has no effect on the 10-11 o'clock (Phase I) and 11-12 o'clock (Phase II) stepping phases, as the states of the sensing pins have been latched earlier.

## Creation of the PRN

An input letter is encrypted by rotating the print head an arbitrary number of steps, resulting in an output letter that is printed onto a paper strip. The arbitrary number of steps is a pseudo-random number (PRN) between 0 and 25 , that is generated by sensing the state of the pins at a certain position on the circumference of six irregularly stepping pin-wheels. If a pin is active, the value of that wheel is added to a running total. This is done by a cylindrical drum - known as the 'cage' - which is in fact a mechanical adder. The cage has 25 horizontally movable bars, each of which has a different set of lugs at one of its long sides. Figure 15 shows a top view of the machine with its cover taken off, in which the cage and the pin-wheels are clearly visible.


Figure 14 - Bar with a lug at the position of pin-wheel 1
The lugs are positioned in such a way that each of them can interact with the sensing arm that determines the state of a pin of a particular pin-wheel, 8 positions above the window. Figure 14 shows a bar with a lug at the position of wheel 1 . The possible positions for lugs that can interact with the sensing arms of wheels $2,3,4$ and 5 are shaded. The lugs for wheel 6 are different, as they are used to invert the bitmask of the calculated sum. When the cage makes one full revolution, all 25 bars are passed by the pin-wheels. The total number of lugs that pass by a particular pin-wheel, determines the value of that wheel. As the lugs are fixed in place - they are not configurable - the value of each wheel (i.e. the sum of the lugs) is fixed as well.


Figure 15 - Top view of the cage with the sideways movable bars

In the SG-41, the lugs on the bars are chosen in such a way that the wheels have the values shown in Figure 16.

- If wheel 1 has an active pin, the print head is rotated by 1 step
- If wheel 2 has an active pin, the print head is rotated by 2 steps
- If wheels 3,4 or 5 , have an active pin, the print head is rotated 4,8 and 10 steps respectively
- If wheel 6 has an active pin, the state of pins 1 to 5 is inverted (i.e. active becomes inactive and vice versa)


## A few examples:

- If the state of the sensed pins of the six pin-wheels is 100110 , the sum is calculated as $1+8+10=19$.
- If the state of the pin-wheels is 100111 , the value is inverted (i.e. 011000 ) and the sum is $2+4=6$.
- If the state of the pin-wheels is $10010 \underline{0}$, the sum is calculated as $1+8=9$.

In the case of these examples, the print head will be advanced by 19,6 or 9 steps respectively.


Figure 16 - Value of each wheel (bitmask)

## Horizontal movement of the bars

The horizontal movement of the bars in the cage of the SG-41 is different from the movement of the bars in a Hagelin cipher machine like the M-209 and BC-38. Instead of two possible positions (in the Hagelin machines) the bars of the SG-41 have thr ee possible positions. Figure 17 shows that each bar has a spring that pulls it to the left. At the left side of the cage is a static cylindrical ramp that engages with the transport lug (i.e. the leftmost lug) of each bar. When rotating the cage, the ramp pushes the bars to the right. Figure 17 shows a bar with an extra lug at the position of wheel 1 , at various angles of rotation of the cage.

As soon as a bar reaches the end of the ramp, it is pulled to the far left by its spring, and the transport lug (i.e. the leftmost lug on each bar) engages with a small cogwheel that can rotate the print head by one position. The sensing arms of the pin-wheels are shaped in such a way that they can interact with the corresponding lugs on the bars. Depending on their state, they may (partially) restrict the movement of the bars. In figure 17, the transport lug of the bar is engaged with the cogwheel.


Figure 17 - Sliding bar shown at different positions whilst the cage revolves

As the lugs on the bars have a lowered and a raised part (and may even be absent), there are three possible positions in which the bar can be locked by a lever. These states are shown in Figure 18. If the lug is absent, or the lever is not engaged, the bar is pulled to the leftmost position (1) and the transport lug overshoots the cogwheel. If the lever catches the raised part of the lug, the bar is locked in the middle position (2) and the transport lug engages with the cogwheel. This is the only situation in which the print head can make a step. If the lever catches the lowered part of the lug, the bar is locked in the rightmost position (3) and the transport lug does not reach the cogwheel.


Figure 18 - Three possible positions of a bar

## Sensing arms

In order to carry the state of the sensed pin of a particular pin-wheel on to the bars in the cage, a complex construction of arms and sickle-shaped levers - pivoting on a common axle - is mounted between the pin-wheels and the cage. This construction is shown in Figure 19. The upper part of the sickle-shaped levers can interact with the bars during phase 1 of the cage rotation. The lower part of the levers is somewhat recessed and can interact with the bars during phase 2 of the cage rotation. If a pinwheel has an active pin, the corresponding sickle-shaped lever is pushed forward by the sensing arm.


Figure 19 - Sickle-shaped sensing arms

## Inactive pin

If the pin of a sensed pin-wheel is inactive (' 0 '), the sickle-shaped lever is not pushed forward. During phase 1 of the cage rotation, the upper part of the lever catches the raised part of the corresponding lug on the current bar. This locks the bar in the middle position (2). During phase 2 of the cage rotation, the lower - recessed - part of the lever does not engage with the lug and allows the bar to be pulled to the leftmost position (1), in which case it overshoots the cogwheel. This means that in this situation, the print head is not rotated.

## Active pin

If the pin of a sensed pin-wheel is active (' 1 '), the sickle-shaped lever is pushed forward. During phase 1 of the cage rotation, the upper part of the lever catches the deeper part of the corresponding lug on the current bar. This locks the bar in the rightmost position (3). During phase 2 of the cage rotation, the lower - recessed - part of the lever catches the raised part of the lug and locks the bar in the middle position (2), coupling the bar with the cogwheel. In this situation, the bar drives the cogwheel, which in turn causes the print head to make a single step.

## Inversion

The rightmost pin-wheel (wheel 6) drives the inversion-function, which is carried out directly by the cage. If wheel 6 has an active pin, the state of the pins of wheels 1 to 5 is inverted. This means that an active pin becomes inactive and vice versa. In this case, the sickle-shaped lever of wheel 6 is pushed forward. During phase 1 of the cage rotation, this has no effect, as the upper part of lever 6 is recessed. In phase 2 of the cage rotation however, the lower part of the lever catches one of the two small lugs at the far right of each bar, which locks the bar in the state it had during phase 1.

Figure 20 shows the four possible situations for wheel 1 only - with an active or inactive pin, with and without inversion. For each situation, phase 1 and phase 2 of the cage rotation are shown. When wheel 6 is inactive, the bar follows the state of the corresponding pin-wheel. This situation is shown in the left half of Figure 20. When wheel 6 is active, the bars are locked in their phase 1 state, by the phase 2 part of the sickle-shaped lever of wheel 6 , that has been pushed forward. This situation is shown in the right half of Figure 20.


Figure 20 - Behaviour when wheel $6=$ ' 0 ' (left) and when wheel $6=' 1$ ' (right)

## Simulation

In order to verify the operation of the SG-41 as explained in the previous chapters, a simple simulation was made in a spreadsheet. It was used to check the stepping mechanism, the order of events and the encryption algorithm, for a finite number of steps. It is the intention of the authors to make this spreadsheet available on the Crypto Museum website in due course. As an example we will show the encryption of a 32 -character plaintext, using a given setting of the machine. First, set the pins as specified in Table 5, in which ' 0 ' means 'left' and '1' means 'right'. Next, set the wheels to the start position highlighted in red:

$$
\begin{array}{llllll}
\text { A } & \text { B } & \text { C } & \text { D } & 01 & 00
\end{array}
$$

Table 6 shows the position of each wheel before encrypting a given 32-character plaintext (PT) - shown in the yellow column. If all goes well, it should produce the specified ciphertext (CT) - shown in the blue column:

PT: SCHLUESSELGERAETVIEREINSWANDERER
CT: IHEPLRETQSDSNDCWHP IVVGLYMHOWSJQS


Table 5 - Pin-settings ( $0=$ left, $1=$ right $)$

Set the wheels to the start position shown in red at line 1 (A B C D 0100 ). This is known as the Grundstellung (basic setting). Enter the first letter of the plaintext $(S)$ and turn the crank by one full revolution. This should yield the first letter of the ciphertext (I). The wheels should now be in the positions specified at line 2 (C E G F 0407 ). Now enter the second letter (C) and turn the crank again. This should yield the second ciphertext letter (H). The window now shows D G H H 05 12. And so on.

Working simulators, that have been verified with these settings, are available for Python [5], Java [6] and BBC BASIC [7].


Table 6 - Wheel positions when encrypting the given plaintext (PT) into ciphertext (CT)

## Conclusions

In this paper we have shown that Schlüsselgerat 41 (SG-41) is based on the design principle of contemporary Hagelin machines like the M-209 and BC-38, but that it contains a number of advanced features which were not available on other machines at the time. We have presented a number of technical details about this machine that have not been available in the public domain before, and we have explained the complex order and timing of the events that take place when a letter is enciphered. In addition we have shown that the common understanding of the wheel stepping principle, namely that the wheels can step in both directions [2], is incorrect. The most important differences with the contemporary Hagelin machines are listed here:

- Irregular stepping of the cipher wheels, with complex timing
- Each wheel has two sensing positions: one for stepping the adjacent wheel and one for rotating the print head
- The sensing position for wheel stepping is re-sensed during the enciphering process
- Improved printer with fixed assignment and (optional) automatic spacing of 5-letter groups
- Inversion of the value generated by wheels 1-5 under control of wheel 6

The principle of irregularly stepping cipher wheels was later introduced by Hagelin on the CX-52 [3].
The inversion principle was used in a modified form in the HELL H-54, which was basically a clone of Hagelin's CX-52. It was produced for the German Army (Bundeswehr) by Hell in Kiel (Germany), under licence from Crypto AG (Hagelin). The inversion principle was known as the 'Hüttenhain feature', and was named after Dr. Erich Hüttenhain who was the head of the Zentralstelle für das Chiffrierwesen (ZfCh) - the German cipher authority. During WWII, Hüttenhain had been the head of the cryptanalysis unit at OKW/Chi - the cipher department of the High Command of the German Wehrmacht (OKW) [4].


Figure 21 - HELL H-54 (left) and Hagelin CX-52 (right)

## About the authors

1. Klaus Kopacz is a self-employed electronics engineer in Stuttgart (Germany). Over the years he has collected a wide range of historical cipher machines, and has helped others to collect, maintain and restore such machines as well. His interests include the technical operation of cipher machines and their historical context. He is particularly known for his Enigma Rebuild arguably the most accurate in its class - and his expertise in the restoration of Enigma and other cipher machines.
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## Links

- Crypto Museum, Schlüsselgerät 41 https://www.cryptomuseum.com/crypto/sg41/

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